Consider the program below, and its control-flow graph.

```
int foo(
                                bb1:
int b0.
                               x = 0
int b1.
                                i = icmp ne b0.0
                               bri, bb2, bb3
int ard3.
int arg4
                                                       bb3:
                                                       i = icmp ne b1.0
int x = 0;
                                                       r = 0
                             bb2:
if (b0) {
                                                       bri, bb4, bb5
                             x = add arg3, arg4
  x = arg3 + arg4;
} else {
  if (b1) {
                                         bb4:
     return 0;
                                         y = add arg3, arg4
                                         r = add x, y
int y = arg3 + arg4;
                                                                 bb5:
return x + y;
                                                                  ret r
```

Figure 1: Example of a program that contains a partial redundancy, and its control-flow graph.

- 1. We say that the program above contains a "partial redundancy". Can you spot the redundant computation in the code above?
- 2. Can you rewrite the program, so that it no longer contains this redundant code?
- 3. Is your optimization performance safe? I mean, is it possible that now, depending on the path that the program takes, it is actually running code that it should not run otherwise?
- 4. We say that an expression is "available" at a program point p if, regardless of the path the program takes to reach p, this expression will be computed before. What does available expressions have to do with redundant code?
- 5. We say that an expression is "very busy (VB)" at a program point p if, regardless of the path the program takes from p till the end of it, this expression will be calculated. What does VB expressions have to do with the performance safety of redundancy elimination?

